### Graph Isomorphism course, Spring 2017

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# 13 Day 13, TuWk7

Homework from last time asked you to show that GI  $\propto_{KARP}$  BIPARTITE GI.

Let  $X \mapsto X'$  be the map on graphs that takes a graph X = (V, E) and returns X', which is found by the addition of |E| new nodes that bisect each edge of X.

**DO 13.1.** Show that  $X \cong Y$  if and only if  $X' \cong Y'$ .

**DO 13.2.** If X is connected and trivalent and e is an edge in X, then  $G = (\operatorname{Aut} X)_e$  is a 2-group.

*Proof.* Write X = (V, E). Consider the Orbit-Stabilizer Theorem:  $|G| = |x^G||G_x|$ .

Let  $T \subseteq V$ . Let  $G_{(T)}$  denote  $(\operatorname{Aut} X)_{(T)}$ , the pointwise stabilizer of T. Suppose that  $v \notin T$  is a neighbor of a vertex in T. If  $T' = T \cup \{v\}$ , by orbit  $v^{G_{(T)}}$  has length 1 or 2. So,  $|G_{(T)}: G_{(T')}|$  is 1 or 2.

Let  $T_0 = e = \{v_1, v_2\}$ , so  $G = G_{T_0}$ . Let  $T_{i+1}$  be built from  $T_i$  by adding a neighbor v of  $T_i$ . By the above argument,  $|G_{T_{i+1}}:G_{T_i}|$  is 1 or 2. Since G is a connected finite graph,  $G_{(T_k)} = 1$  for some k. We have found that |G| is a power of 2.

### 13.1 Reduction of Trivalent GI to SI of 2-groups

We start with a set of vertices T. Denote by N(T) the set of neighbors of T. We construct a layered graph:  $L_0 = T$ ,  $L_1 = N(T) \setminus T$ , and  $L_k = N(L_{k-1}) \setminus L_{k-1}$  for all k.

Two vertices in  $L_{k+1}$  are **twins** if their neighbors in  $L_k$  are identical.

Assume  $G_{k-1} = \operatorname{Aut}(X_{k-1})$  is known. Suppose we want to compute  $G_k$ . Let  $\varphi: G_k \to G_{k-1}$  be the restriction. This is a homomorphism because  $X_{k-1}$  is invariant under  $G_k$ . Notice that  $\operatorname{Ker}(\varphi)$  is the direct product of symmetric groups on twin equivalence classes. We are interested in finding  $\operatorname{Img}(\varphi) \leq G_{k-1}$ .

**DO 13.3.** Show that  $G_k = \langle \text{lifting of each generator of Img } \varphi, \text{ generators of Ker} \rangle$ .

For each  $L_k$ , build the hypergraph on  $L_k$  given by the set of neighbors of vertices in  $L_{k+1}$ .

**DO 13.4.** If  $L_{k-1}$  has m vertices, then reduce finding  $\operatorname{Img}(\varphi)$  to SI for strings of length  $m + {m \choose 2} + {m \choose 3}$ .

#### 13.2 GI for bounded degree graphs is polynomial time

**DO 13.5.** Suppose that X is connected and d-valent. Show that  $G = (\operatorname{Aut} X)_e$  is in  $\Gamma_{d-1}$ .

We give two definitions of  $\Gamma_d$ , a class of finite groups.

**Definition 13.6.** A group G is in  $\Gamma_d$  if there exists a subgroup chain of G

$$G = G_0 \ge G_1 \ge \cdots \ge G_m = 1$$
,

such that  $|G_{i-1}:G_i| \leq d$ .

**Definition 13.7.** A group G is in  $\Gamma_d$  if every composition factor embeds in  $S_d$  (isomorphic to a subgroup of  $S_d$ ).

**HW 13.8.** Show that the two definitions of  $\Gamma_d$  are equivalent.

Notice that if  $G \in \Gamma_d$ , then  $\Theta(G) \leq d$ . The following Theorem follows as a corollary from the BCP-Theorem. It was shown previous to the BCP-Theorem by Luks via a different trick.

**Theorem 13.9** (Luks (1980)). GI for graphs of degree  $\leq d$  is in polynomial time.

## 13.3 Tournaments and Petersen Graph

A **tournament** is an orientation of the complete graph – a directed graph on V where, for all  $u, v \in V$  with  $u \neq v$ , either  $(u, v) \in E$  or  $(v, u) \in E$ .

**HW 13.10.** If X is a tournament, then Aut(X) is solvable.

Challenge 13.11 (B-Luks (1983)). Isomorphism of tournaments is decidable in  $n^{O(\log n)}$ .

**HW 13.12.** The automorphism group of the Petersen graph is isomorphic to  $S_5$ .

Central to the new GI algorithm is the "Unaffected stabilizers lemma." Look it up on arXiv.