

## Graph Isomorphism course, Spring 2017

Instructor: László Babai

Notes by Angela Wu and the instructor

Tuesday, May 23, 2017

## 17 Day 17, TuWk9

### 17.1 Symmetry defect of regular graphs

**Proposition 17.1.** *If  $X$  is a nontrivial regular graph, then its symmetry defect is  $\geq 1/2$ .*

Recall that we say  $u, v \in V$  are **twins** if the transposition  $(u, v) \in \text{Aut}(X)$ . Also,  $\text{defect}(X) = 1 - |T|/n$ , where  $T$  is the largest equivalence class of  $V(X)$  under the equivalence relation of “twin or equal.”

Notice that  $\text{defect}(K_n) = 0$ .

**DO 17.2.**  $\text{defect}(X) = \text{defect}(\bar{X})$ , because  $\text{Aut}(X) = \text{Aut}(\bar{X})$ .

**Lemma 17.3.** *If  $X$  is  $k$ -regular and nonempty ( $k \geq 1$ ), then all independent sets have size  $\leq n/2$ .*

**Lemma 17.4.** *If  $T$  is a set of twins in a graph  $X$ , then the induced subgraph  $X[T]$  is empty or complete ( $T$  is either independent or a clique).*

Indeed,  $\text{Aut}(X[T]) = \text{Sym}(T)$ .

*Proof of Prop. 17.1.* Let  $T$  be a twin equivalence class. We need to show that  $|T| \leq n/2$ . Indeed, if  $T$  is an independent set then  $|T| \leq n/2$  by Lemma 17.3. If  $T$  is a clique then  $T$  is an independent set in the complement of  $X$  and therefore again  $|T| \leq n/2$ .  $\square$

Note: this proof is simpler than the one given in class.

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For the remaining exercises, read the appropriate sections of the instructor’s paper (accessible from the course home page). Below we refer to that paper as [GIQ]. (Why?)

### 17.2 Naive refinement

One round of the naive vertex refinement process takes as input a vertex-colored graph  $X = (V, E, f)$  where  $f : V \rightarrow \{\text{colors}\}$  is the coloring, and outputs a refined coloring  $g$  defined as follows:

$g(v) = (f(v); d_i(v) \mid i \in \{\text{colors}\})$  where  $d_i(v)$  denotes the number of neighbors of  $v$  of color  $i$ . The naive refinement process iterated this refinement step until the color partition stabilizes (no number of colors does not increase).

**DO 17.5.** A coloring is stable under naive refinement if and only if the color classes form an equitable partition.

(See [GIQ], Sec. 2.4.2 for the definition of equitable partitions.)

**DO 17.6** (Abe Mowshowitz  $\approx$  1970). If the characteristic polynomial of the adjacency matrix of a graph is irreducible over the rationals then naive refinement completely splits the graph.

Hint: show that under this assumption, only the discrete partition is equitable.

### 17.3 Coherent configurations (CCs)

Read [GIQ], Section 3 for the basics of the combinatorial theory of CCs. (CCs are called “classical CCs” in the paper.) Let  $G \leq \text{Sym}(V)$ . Denote by  $\mathfrak{X}(G)$  the orbital configuration of  $G$  (given by the orbits of  $G$  on  $V \times V$ ).

**DO 17.7.** Show that the orbitals of a group  $G$  form a coherent configuration.

**DO 17.8.**  $\mathfrak{X}(G)$  is homogeneous if and only if  $G$  is transitive.

**HW 17.9.**  $\mathfrak{X}(G)$  is primitive if and only if  $G$  is primitive.

**DO 17.10.** Show that the Johnson scheme  $\mathfrak{J}(k, t)$  is a primitive coherent configuration. — Give two proofs: one combinatorial (directly from the definition), the other using the fact that the group is  $S_k^{(t)}$  is primitive. (Hint for the second proof: the Johnson schemes are Schurian.)

**DO 17.11.**  $\text{Aut}(\mathfrak{X}(G)) \geq G$ .

**HW 17.12.** (Typo fixed and additional explanation added at 8pm on 5-20.) Find an infinite family of groups  $G \leq S_n$  (for infinitely many values of  $n$ ) such that  $|G|$  is small (grows slowly as a function of  $n$ ) but  $|\text{Aut}(\mathfrak{X}G)|$  is very large (grows fast as a function on  $n$ ). State conceptually what “small” and “very large” mean in terms of commonly used rates of growth (as functions of  $n$ ).

**HW 17.13.** Prove that  $c(x, y)$  is aware of  $c(x)$  and  $c(y)$ .

**DO 17.14.** Each constituent of a CC is either biregular digraph or a semiregular bipartite graph.

The *directed distance* from vertex  $u$  to vertex  $v$  in a digraph is the length of the shortest directed path from  $u$  to  $v$ . The *directed diameter* of a digraph is the largest directed distance among pairs of vertices. (This quantity is infinite if the graph is not strongly connected.)

**HW 17.15** (Due Tuesday 5/30). All connected components of a homogeneous constituent of a CC have the same directed diameter.

(They also have the same undirected diameter, but you don’t need to prove this.)

**DO 17.16** (Useful for earlier exercises). Let us fix a sequence  $c_1, \dots, c_t$  of (not necessarily distinct) colors. The number of walks of length  $t$  of color sequence  $(c_1, \dots, c_t)$  from  $x$  to  $y$  is determined by  $c(x, y)$  and does not depend on the specific choice of  $x$  and  $y$ .