

PROBLEM SESSION | 2023-11-17

1

14.65 6-color theorem: \forall planar graph is 6-colorable

(1) If G is planar and $H \subseteq G$ then H is planar

(2) If G is planar, $n \geq 1$, then \exists vertex of $\deg \leq 5$

Proof induction on n

$n = 1$ ✓

$n \geq 2$, assume true for $n-1$ vertices \leftarrow IH

G has n vertices, pick a vertex x of $\deg \leq 5$

$G - x$ 6-colorable by IH

Claim Any 6-coloring of $G - x$ can be extended
to a 6-coloring of G .

Pf Pick a 6-coloring of $G - x$. Now color x by
a color not used by its neighbors. ✓

14.89

$$\binom{n}{5} = \frac{n(n-1)\cdots(n-4)}{5!} = \frac{N}{5!}$$

Don't

$$\binom{n}{5} = \frac{n!}{5!(n-5)!}$$

true but ugly

$$\frac{N}{n^5} = \frac{n(n-1)\cdots(n-4)}{n^5} = \left(1 - \frac{1}{n}\right)\left(1 - \frac{2}{n}\right)\left(1 - \frac{3}{n}\right)\left(1 - \frac{4}{n}\right) \rightarrow 1$$

$\downarrow \quad \downarrow \quad \downarrow \quad \downarrow$
 $1 \quad 1 \quad 1 \quad 1$

$$\therefore N \sim n^5$$

$$\therefore \binom{n}{5} = \frac{N}{5!} \sim \frac{n^5}{5!} = \frac{n^5}{120}$$

alternative proof:

N is a poly of deg 5 in n
 with leading term n^5
 $\therefore N \sim n^5$ by a general result

$$\left(1 - \frac{1}{n}\right)^n = \left(1 - \frac{1}{n}\right)\left(1 - \frac{1}{n}\right)\cdots\left(1 - \frac{1}{n}\right) \rightarrow e$$

$\downarrow \quad \downarrow \quad \downarrow$
 $1 \quad 1 \quad 1$

14.96 $f(x) = ax^k + bx^{k-1} + \dots \quad ab \neq 0$

[3]

then $f(n) \sim an^k$

14.96 $\frac{f(n)}{n^k} = a + \frac{b}{n} + \frac{\dots}{n^2} + \dots \rightarrow a$

$\downarrow \quad \downarrow \quad \downarrow$
 $0 \quad 0 \quad 0$

$\frac{f(n)}{a \cdot n^k} \rightarrow 1 \quad f(n) \sim an^k$

14.97 $\frac{f(x)}{g(x)} = \frac{ax^k + \dots}{cx^l + \dots}$

$\frac{f(n)}{g(n)} \sim \frac{an^k}{cn^l} = \underline{\underline{\frac{a}{c} \cdot n^{k-l}}}$

14.101

4

$$\sqrt{n^2+1} - n = \frac{1}{\sqrt{n^2+1} + n} \sim \frac{1}{2n}$$

$$\sqrt{n^2+1} \sim n$$

$$\sqrt{n^2+1} + n \sim 2n$$

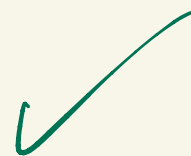
Divide by n :

$$\sqrt{1 + \frac{1}{n^2}} + 1 \rightarrow 2$$

↓
0

$$\frac{\sqrt{n^2+1} + n}{2n} \rightarrow 1$$

∴



14.99

(5)

$$\ln\left(1 + \frac{1}{n}\right) \sim \frac{1}{n}$$

~~Apply L'Hospital~~

~~to $\frac{\ln\left(1 + \frac{1}{n}\right)}{\frac{1}{n}}$~~

Lemma

$$\lim_{x \rightarrow 0} \frac{\ln(1+x)}{x} =$$

$$= \lim_{x \rightarrow 0} \frac{\frac{1}{1+x}}{1} = \lim_{x \rightarrow 0} \frac{1}{1+x} = 1$$

then follows by setting $x = \frac{1}{n}$

Alternative proof: ^{of Lemma} instead of L'Hospital,
use def of derivative:

Let $f(x) = \ln x$

$$f'(1) = \lim_{h \rightarrow 0} \frac{\ln(1+h) - \ln(1)}{h}$$

$$f'(a) = \lim_{h \rightarrow 0} \frac{f(a+h) - f(a)}{h}$$

$$\therefore \lim_{h \rightarrow 0} \frac{\ln(1+h)}{h} = \ln'(1) =$$

$$\left. \frac{1}{t} \right|_{t=1} = 1 \quad \checkmark$$

$$14.99 \quad \sin\left(\frac{1}{n}\right) \sim \frac{1}{n}$$

LEMMA

$$\lim_{x \rightarrow 0} \frac{\sin x}{x} = 1$$

b/c this limit is $(\sin)'(0) = \cos 0 = 1 \quad \checkmark$

(also works by L'Hôpital)

$$14.105 \quad a_n \sim b_n \not\Rightarrow a_n^n \sim b_n^n$$

Counterexample: $a_n = 1, \quad b_n = 1 + \frac{1}{n}$

$$a_n^n = 1, \quad b_n^n = \left(1 + \frac{1}{n}\right)^n \rightarrow e$$

14.131

7

$$a_n, b_n > 1$$

$$a_n \sim b_n \not\Rightarrow \ln a_n \sim \ln b_n$$

Pf : need counterexample : pair of sequences that satisfies $a_n \sim b_n$ but $\ln a_n \not\sim \ln b_n$

$$\left. \begin{array}{l} a_n := e^{1/n} \rightarrow 1 \\ b_n := e^{2/n} \rightarrow 1 \end{array} \right\} \Rightarrow a_n \sim b_n$$

$$\ln a_n = \frac{1}{n}$$

$$\ln b_n = \frac{2}{n}$$

$$\text{quotient} = \frac{1}{2} \rightarrow \frac{1}{2} \quad \checkmark$$

Another example: $a_n = 1 + \frac{1}{n} \rightarrow 1$
 $b_n = 1 + \frac{2}{n} \rightarrow 1 \quad \therefore a_n \sim b_n$

$$\left. \begin{array}{l} \ln a_n \sim \frac{1}{n} \\ \ln b_n \sim \frac{2}{n} \end{array} \right\} \Rightarrow \frac{\ln a_n}{\ln b_n} \sim \frac{1}{2}, \text{ i.e. } \frac{\ln a_n}{\ln b_n} \rightarrow \frac{1}{2}$$

14.37 Greedy coloring uses $\leq 1 + \Delta$ colors

$$\Delta = \deg_{\max}$$

Pf induction on n . True if $n=1$. Now $n \geq 2$

IH true for $n-1$.

G has n vertices

\nearrow
input, given to us by an adversary

pick any $x \in V(G)$

$$\Delta(G-x) \leq \Delta(G)$$

so by IH $G-x$ is colorable by $\Delta(G-x)$ colors: by $\Delta(G)$ colors

Pick a Δ -coloring of $G-x$

Claim This coloring extends to G

Pf: color x by a color not used by any neighbor. ✓

(8)

1 even

9

14.39 Find bipartite graph s.t.

greedy col. uses $n/2$ colors

Pf $V = [n]$

for $k = 1, \dots, \frac{n}{2}$

$2k-1$ is adjant to $2j$ for all $j \neq k$

Claim $\chi \subseteq$ denotes

greedy coloring then

$(\forall k) (c(2k-1) = c(2k) = k)$

Pf by induction k

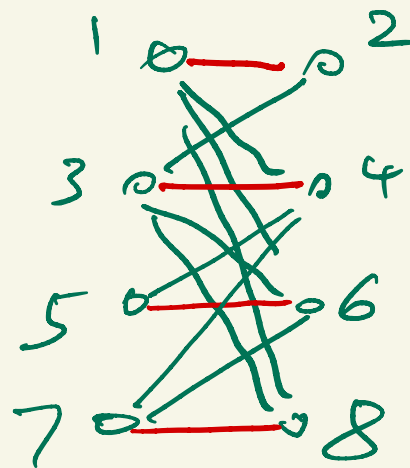
Base: $k=1$ ✓

IH: true for all $k' < k$

If for k : $\forall x$ $2k-1$ is adj.

to $2j$ for $j \leq k-1$ - these have all colors $1, \dots, k-1$ so
 $c(2k-1) = k \leftarrow$ forced

Same
for
 $\forall x$ $2k$ ✓



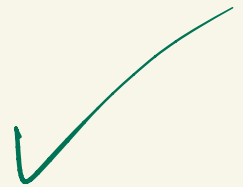
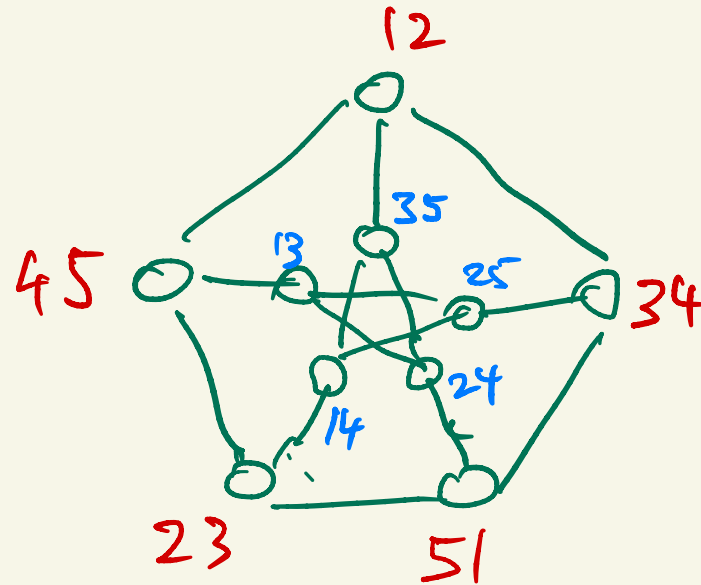
$r \geq 2s$ Kneser's graph

[10]

$$V = \mathcal{V}(K_n(r, s)) = \binom{[r]}{s}$$

$$A, B \in V \quad A \sim B \text{ if } A \cap B = \emptyset$$

$$K_n(5, 2) \cong \text{Petersen's}$$



14.49

 $K_n(r, s)$ is regular of deg $\binom{r-s}{s}$

(A)

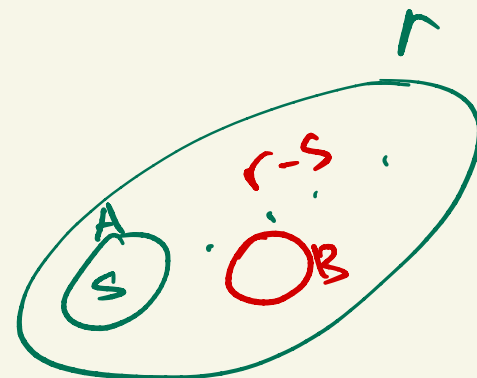
$$|A| = s$$

$$A \subseteq [r]$$

count $B \subseteq [r]$

$$|B| = s$$

$$A \cap B = \emptyset$$



$$14.53 \quad \chi(K_n(r, s)) \leq r - 2s + 2$$

(12)

vertices of each color must be an independent set
 indep. set in K_n corresponds to
 s -subsets that pairwise intersect

C_1 : all s -subsets containing "1"

C_2 : " " - not containing "1" but containing "2"

\vdots

C_k : " " - not containing $1, \dots, k-1$ - " " - "k"

for $k=1, \dots, r-2s+1$

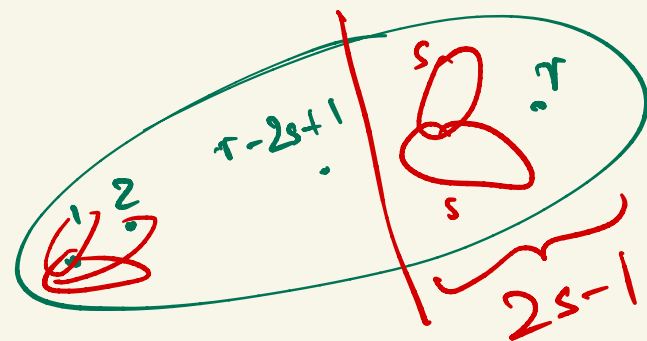
remaining :

s -subsets of $\{r-2s+2, \dots, r\}$
 $2s-1$ elements



pairwise intersect

→ get 1 color



$r-2s+1$ colors 1 color

~~If $a_n \sim b_n$ then $\lim a_n = \lim b_n$~~

$$15.47 \quad (\ln n)^{100} = o(n)$$

$$\text{NTS} \quad \frac{(\ln n)^{100}}{n} \rightarrow 0$$

$$\left[\text{i.e. } \frac{\ln n}{n^{1/100}} \rightarrow 0 \right]$$

$$15.45 \quad \underline{\text{DO}} \quad (\forall c > 0) \quad \lim_{x \rightarrow \infty} \frac{\ln x}{x^c} = 0$$

\Downarrow

$$15.47 : \quad c := \frac{1}{100}$$

bootstrapping

Pf: L'Hôpital once

$$\frac{1/x}{c x^{c-1}} = \frac{1}{c} \cdot \frac{1}{x^c} \rightarrow 0$$

✓

Actually $15.43 \Rightarrow 15.45 \quad \forall c > 0$

$$\frac{\ln x}{x} \rightarrow 0 \quad x \rightarrow \infty$$

\Rightarrow

$$\frac{\ln y}{y^c} \rightarrow 0$$

$$\frac{\ln y}{y^c} = \frac{1}{c} \cdot \frac{\ln x}{x}$$

use 15.43 with $x = y^c$

$$\ln x = c \cdot \ln y$$

$$\ln y = \frac{1}{c} \ln x$$

"exponential growth beats poly. growth"

$$f(x) = e^{x^c}$$

$$g(x) = x^c$$

then $\frac{g(x)}{f(x)} \rightarrow 0$

example: $e^{\sqrt{x}}$

If case $c=1$ ~~⊗~~

$$\frac{x^c}{e^x} \stackrel{?}{\rightarrow} 0$$

$y := e^x$
 $x = \ln y$

$$\frac{(\ln y)^c}{y} \rightarrow 0$$

✓

Case general $c > 0$

$$\frac{x^c}{e^{x^c}} \stackrel{?}{\rightarrow} 0$$

$z := x^c$
 $x = z^{1/c}$

$$\frac{z^{c/c}}{e^z} \rightarrow 0 \quad b/c \quad \otimes$$

15.15 Given $(b_n) \exists (a_n)$ s.t. $a_n = o(b_n)$

$\iff b_n$ is eventually nonzero

Pf. \implies ASSN $a_n = o(b_n)$, i.e. $\frac{a_n}{b_n} \rightarrow 0$
 NTS b_n event. nonz.

Pf by contradiction: $b_n = 0$ inf. often $\therefore \frac{a_n}{b_n}$ is undefined
 inf. often $\therefore \nexists \lim \frac{a_n}{b_n}$

\Leftarrow ASSN b_n event. nonzero:
 $(\exists n_0)(\forall n > n_0)(b_n \neq 0)$

NEED (a_n) s.t. $\frac{a_n}{b_n} \rightarrow 0$ let $a_n = 0$

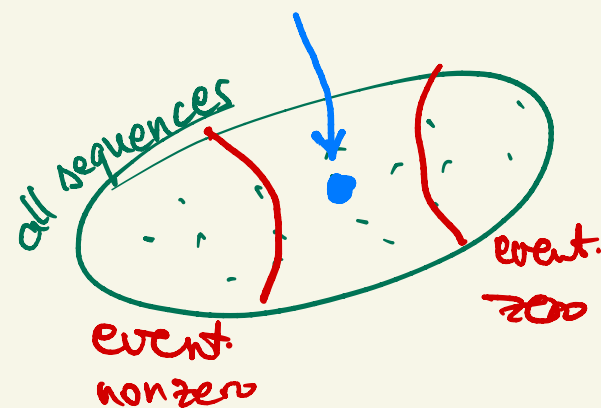
$\rightarrow b_n = 0$ infinitely often

we find sequence that is
neither

010101

$\neg (b_n \text{ is eventually nonzero})$

is NOT " b_n is eventually zero"



ASSNs

14.129

$$a_n, b_n > 1$$

$$a_n \sim b_n$$

a_n bdd away from 1

DC b_n bdd away from 1

Assn $(\exists \varepsilon > 0)(\exists n_0)(\forall n)(n > n_0 \Rightarrow a_n > 1 + \varepsilon)$

NTS $(\exists \delta > 0)(\exists n_1)(\forall n)(n > n_1 \Rightarrow \underline{b_n > 1 + \delta})$

Claim ~~$n_1 = n_0$~~ , any $0 < \delta < \varepsilon$ works

Pf Want $b_n > \frac{1+\delta}{1+\varepsilon} a_n$ this will suffice if $n > n_0$
 \Downarrow
 $a_n > 1 + \varepsilon$
 \Leftarrow
 $b_n > 1 + \delta$

$$\frac{1+\delta}{1+\varepsilon} =: 1 - \gamma \quad \gamma > 0$$

$$a_n \sim b_n \Rightarrow (\exists n_2)(\forall n)(n > n_2 \Rightarrow \frac{b_n}{a_n} > 1 - \gamma)$$

$$n_1 := \max(n_0, n_2)$$