

HONORS

2024-01-12

ALGORITHMS

1

input $a, b \in \mathbb{N}_0 = \{0, 1, 2, \dots\}$

output a^b

bitlength $\approx \log_2 a + \log_2 b$
of input

of output $b \cdot \log_2 a$ $a \geq 2$

exponential

bitlength of n
is $1 + \lfloor \log_2 n \rfloor$
 $\boxed{ID} = \lceil \log_2 (n+1) \rceil$

input a, b, m

output $(a^b \bmod m)$

DEF divisibility $d \mid a$ if $(\exists x)(a = dx)$

↙

in particular, $0 \mid 0$

Congruence

DEF $a \equiv b \pmod{m}$ if $m \mid a - b$

DEF assume $m \neq 0$

$r = (a \bmod m)$ least non-neg. remainder
of division by m

DO $r = (a \bmod m) \iff \begin{cases} 0 \leq r \leq |m|-1 \\ r \equiv a \pmod{m} \end{cases}$

task

given $a, b, m \in \mathbb{Z}$ compute $(a^b \bmod m)$
 $m \geq 1$

TASK Given $a, b, m \in \mathbb{N}_0$ compute $(a^b \bmod m)$ 3
 $m \geq 1$

DO Computing $(x \bmod m)$ can be done
in quadratic time = #bit. operations
"long division"

compute a^b , reduce mod m X exp. time

$x := 1$

$b \geq 1$

for $k = 1$ to b

$x := ax$

return x

returns a^b

improved:

$x := 1$

for $k = 1$ to b

$x := (ax \bmod m)$

rounds = b : exponential in $\log b$

we never have to
deal w numbers $\geq m^2$
($2 \log m$ bits)

REPEATED SQUARING

(4)

$$(a^{32} \bmod m)$$

$x := a$
for $i = 1$ to 5
 $x := (x^2 \bmod m)$

$$a^{39} = a^{32} \cdot a^4 \cdot a^2 \cdot a^1$$

a
 a^2
 a^4
 a^8
 a^{16}
 a^{32}

repeated squaring

done with $2 \cdot \log b$ modular multiplications

\therefore polynomial time

Input : a, b, m ← parameters
auxiliary Variables
 A, B, X ← accumulator

of the algorithm
 they don't change during execution of alg

$$a^b = a \cdot a^{b-1}$$

$$a^{2k} = (a^k)^2$$

initialize $A := (a \bmod m)$
 $B := b$
 $X := 1$

while $B \geq 1$ •
 if B is odd then $B := B - 1, X := (A \cdot X \bmod m)$
 else $B := B/2, A := (A^2 \bmod m)$

return X

loop invariant : statement (Y/N)
 about the configuration (evaluation of variables)
 the if true when we enter an execution of the loop
 then true on exit

initialize $A := (a \bmod m)$
 $B := b$
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6

while $B \geq 1$ •
 if B is odd then $B := B - 1, X := (A \cdot X \bmod m)$
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return X

loop invariant : statement (Y/N)
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variables)
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the loop
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$$X \cdot A^B \equiv a^b \pmod{m}$$

true at beginning

→ by induction: true in each step

⇒ true at the end: $B = 0$ ∴ $X \equiv a^b \pmod{m}$

LOOP INVARIANT

7

loop: while P do T

predicate over a set A

$$\underline{P}: A \rightarrow \{0,1\}$$

\mathcal{C} : configuration space

$$\underline{P}: \mathcal{C} \rightarrow \{0,1\}$$

predicate

$$T: \mathcal{C} \rightarrow \mathcal{C}$$

transformation

$$R: \mathcal{C} \rightarrow \{0,1\} \text{ predicate over } \mathcal{C}$$

DEF R is a loop invariant if

$$(\forall X \in \mathcal{C}) (P(X) \wedge R(X) \rightarrow R(T(X)))$$