

HONORS

2024-01-19

ALGORITHMS

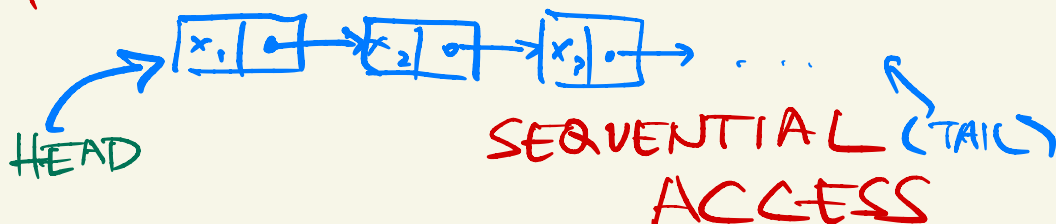
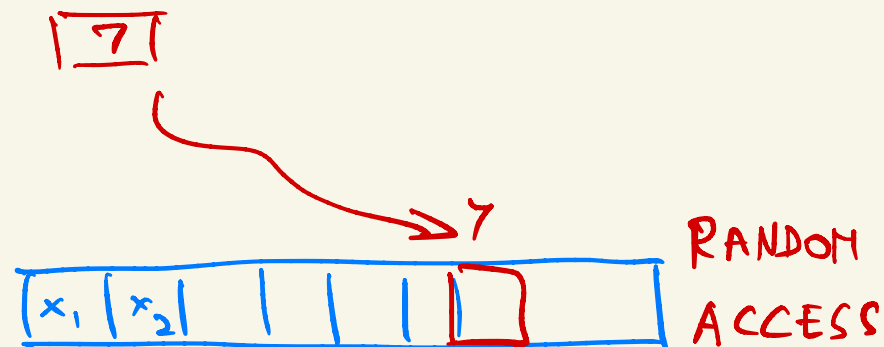
Data structures

LIST

basic types:

ARRAY

LINKED LIST



If SORTED : ARRAY

FIND

binary search

efficient: $O(\log n)$

LINKED LIST

sequential

inefficient: $O(n)$

INSERT/
DELETE

requires shifting
inefficient: $O(n)$

INCREASE/DECREASE
KEY

requires re-sorting

update links
efficient: $O(1)$ } assuming
address
found

COUNTING SORT

Suppose all keys $key(i) \in [m]$

n items

SORT in $O(n)$ operations

$$A = [key(1), \dots, key(n)]$$

$$B = [1 \dots m]$$

$$A = [3, 5, 3, 2]$$

$$B = [1, 2, 3, 4, 5, 6]$$

DO: pseudocode

single pass through A n

" " B m

total $O(n+m)$

B	1	2	3	4	5	6
B*	0	1	2	0	1	0

$$C = [2, 3, 3, 5]$$

linear time if $m = O(n)$

m-way branching

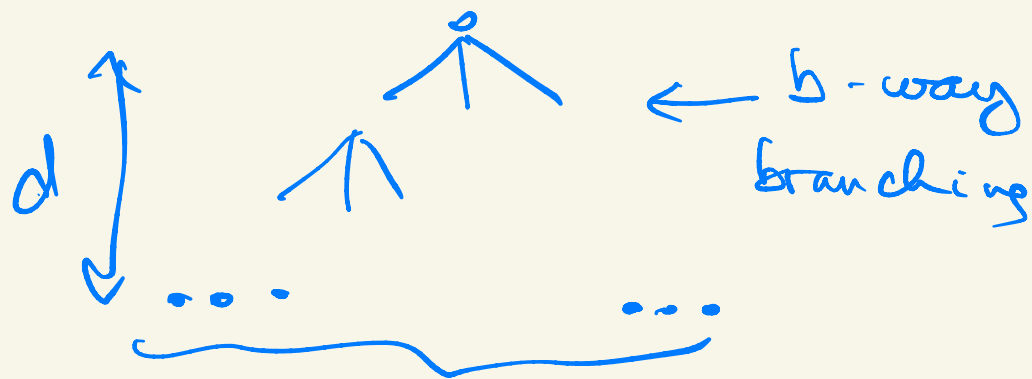
$$d \geq \log_m(n!) = \frac{\log(n!)}{\log m}$$

Case $m = n$

$$d \geq \log_n(n!) = \frac{\log(n!)}{\log n} \sim \frac{n \ln n}{\ln n} = n$$

info thg lower bound

3



n outcomes all distinguished

$$\Rightarrow n \leq b^d$$

$$d \geq \log_b n$$

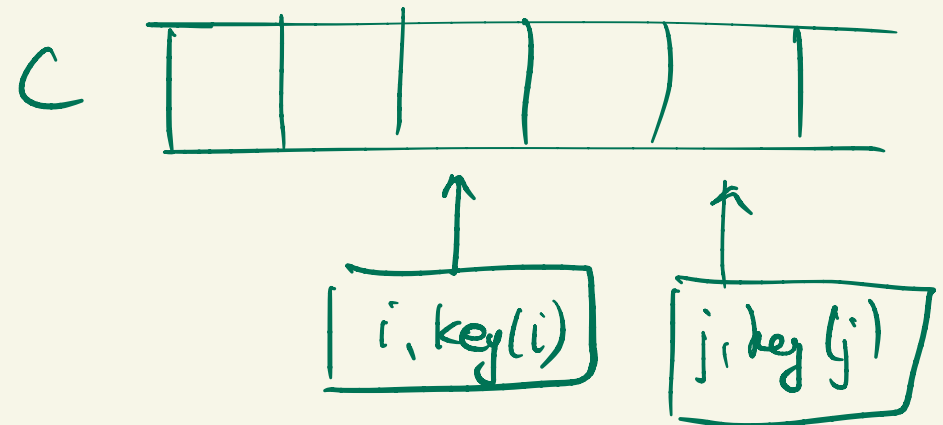
			2		
			3		
			2		

ID
ID

↑
key

Sort array
by column k

4



then $\text{key}(i) \leq \text{key}(j)$
 and if $\text{key}(i) = \text{key}(j)$
 then $i \leq j$

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(a_{ij})

r_1	a_{11}	a_{12}			a_{1m}
r_2	a_{21}	a_{22}			a_{2m}
r_n					

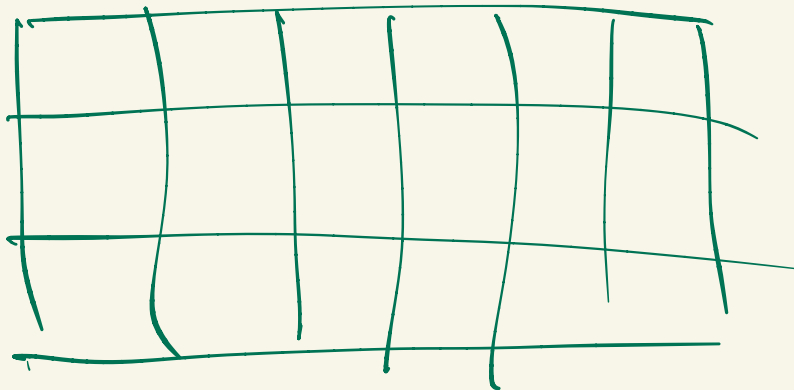
Sort rows lexicographically (dictionary order)

$$r_i < r_j \quad \text{if } (\exists k) (a_{i1} = a_{j1} \dots a_{i,k-1} = a_{j,k-1} \text{ and } a_{ik} < a_{jk})$$

RADIX SORT (6)

(a_{ij}) $i = 1 \dots n$
 $j = 1 \dots q$

$a_{ij} \in [m]$



$n \times q$
array

for $j = q$ downto 1

counting sort by column j

EX verify correctness

COST $O((n+m)q)$

Output:
permutation

$f: [n] \rightarrow [n]$

i.e. bijection
of $[n]$ to itself

for $m = O(1)$ we get $O(nq) \leftarrow$ linear

EX. Sort strings of variable lengths ^{over [m]} in linear time, assuming $m = \text{const.}$

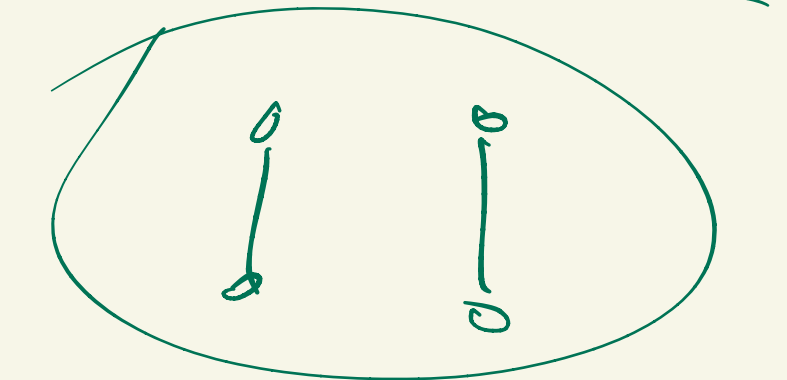
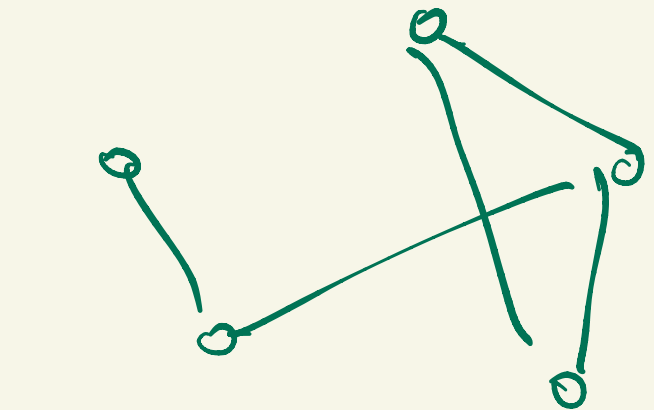
GRAPHS

(nodes, links)

(V, E)

vertices
vertex

edges



Given n vertices, # graphs
on this set of
vertices?

8

edge $\{i, j\}$
 $= \{j, i\}$



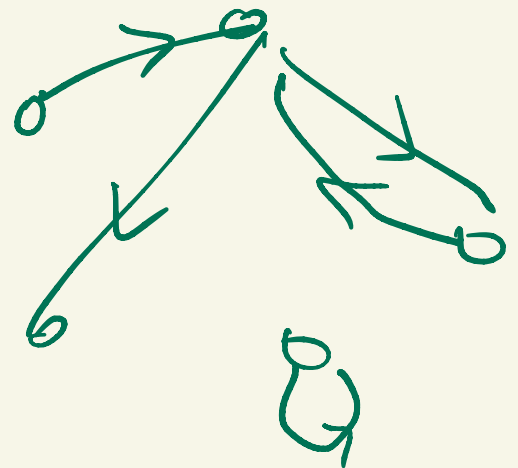
$2^{\binom{n}{2}}$

Digraphs (directed graphs)

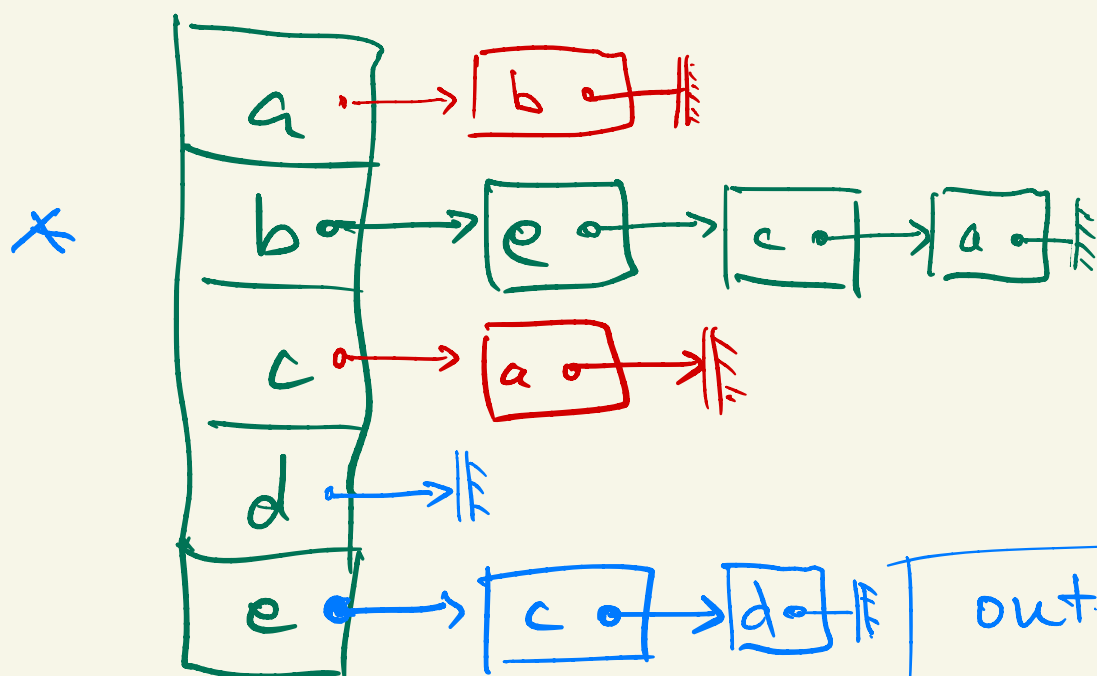
edge (i, j) $i \rightarrow j$

n^2 possible edges

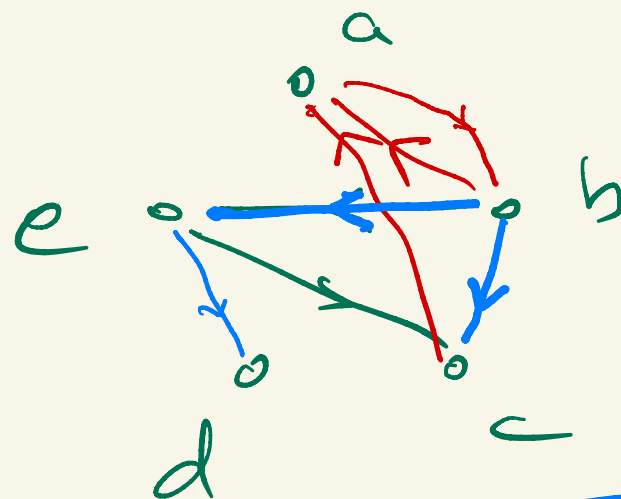
2^{n^2}



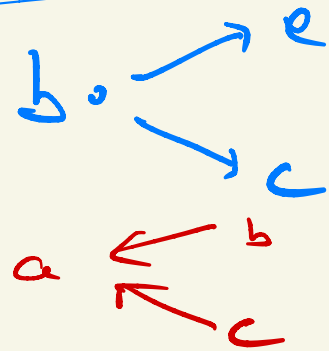
Standard representation: array of linked lists



↑
array of vertices

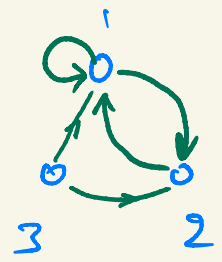


out-neighbors
in-neighbors



$i \rightarrow j$ (i, j) are adjacent

this is the adjacency relation



adjacency matrix $a_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E \\ 0 & \text{o/w} \end{cases}$

$n \times n$ $(0,1)$ -matrix

1	1	0
1	0	0
1	1	0

Size: n^2

linked-list rep: size $O(n+m)$

$m = \# \text{ edges}$

EX. in linear time

- sort all adj. lists
- remove repetitions
- reverse every edge