

HONORS ALGORITHMS

2024-02-09

1

NEW HANDOUT: WALK-TO-PATH

Link from HW page, 15.10

POLYNOMIAL IDENTITY LEMMA

(2)

$$f \in F[x_1, \dots, x_n] \quad F \cdot \text{field}$$

$$H \subseteq F$$

If $f \neq 0$ (formally zero) (all coeff's zero)

$$\text{then } P(f(a_1, \dots, a_n) = 0) \leq \frac{d}{|H|} \quad \leftarrow \text{degree}$$

$\nwarrow$ finite

Proof: induction on n

$$n=1 \quad P(f(a)=0) \leq \frac{d}{|H|}$$

$a \in H$

b/c f has at most d roots in F

now assume $n \geq 2$

$$\deg f = d$$

(3)

I.H. True for $n' < n$ variables

$$A = \{ \underline{a} \in H^n \mid f(\underline{a}) = 0 \} \quad \underline{a} = (a_1, \dots, a_n)$$

$$\underline{\text{NTS}} \quad |A| \leq d \cdot h^{n-1} = \frac{d}{h} \cdot h^n \quad h = |H|$$

$$f(x_1, \dots, x_n) = g_0(\underline{x}) + g_1(\underline{x}) \cdot x_n + \dots + \underline{g_k(\underline{x})} \cdot x_n^k$$

$$\underline{x} = (x_1, \dots, x_{n-1})$$

$$g_k \neq 0$$

$$B = \{ \underline{b} \in H^{n-1} \mid g_k(\underline{b}) = 0 \}$$

$$\deg g_k \leq d - k$$

$$\text{I.H.} \Rightarrow |B| \leq (d - k) \cdot h^{n-2}$$

$$f(x_1 \dots x_n) = g_0(\underline{x}) + g_1(\underline{x}) \cdot x_n + \dots + \underline{g_k(\underline{x})} \cdot x_n^k$$

$$\underline{x} = (x_1 \dots x_{n-1})$$

$$g_k \neq 0$$

$$B = \{ \underline{b} \in H^{n-1} \mid g_k(\underline{b}) = 0 \}$$

$$\deg g_k \leq d-k$$

$$\text{I.H.} \Rightarrow |B| \leq (d-k) \cdot l^{n-2} \quad \otimes$$

$$\text{If } \underline{b} \notin B \text{ then } g_k(\underline{b}) \neq 0$$

$$C(\underline{b}) = \{ a \in H \mid f(\underline{b}, a) = 0 \} \quad |C(\underline{b})| \leq k$$

$$g_k(\underline{b}) \neq 0 \quad \therefore f(\underline{b}, x_n) \text{ has } \deg = k \quad \swarrow$$

$$\# \{ (\underline{b}, a) \mid \underline{b} \in H^{n-1} \setminus B \wedge f(\underline{b}, a) = 0 \} \leq \underline{l^{n-1} \cdot k}$$

$$\# \{ (\underline{b}, a) \mid \underline{b} \in B \wedge \dots \} \leq |B| \cdot l \leq \underline{l^{n-1} \cdot (d-k)} \quad \otimes$$

$$\therefore \# \{ (\underline{b}, a) \mid (\underline{b}, a) \in H^n, f(\underline{b}, a) = 0 \} \leq \underline{l^{n-1} \cdot d} \quad \checkmark$$

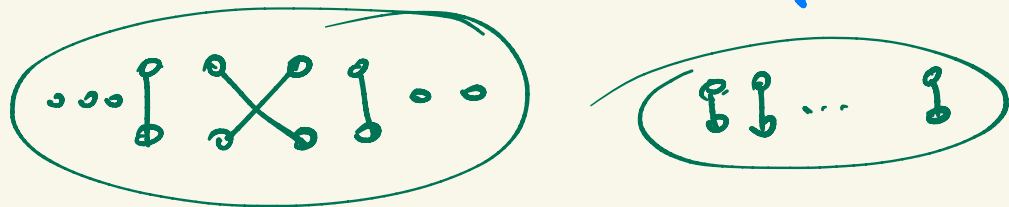
This bound is tight for every field F

and every (n, d, h) s.t. $d \leq h \leq |F|$, $\forall H \subseteq F$
 $|H| = h$

i.e. $\exists f \in F[x_1, \dots, x_n]$, $\forall H \subseteq F$, $|H| = h$
s.t. $P_{a \in H^n} (f(a) = 0) = \frac{d}{h}$

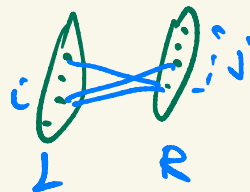
Solution: $f(x_1, \dots, x_n) = \prod_{\alpha \in D} (x_1 - \alpha)$ where $D \subseteq H$
 $|D| = d$

Perfect matching in graph G : matching of size $\frac{n}{2}$

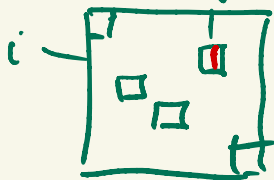


Bipartite graph with $2n$ vertices, n on each side
bipartite adjacency matrix: $n \times n$

$$a_{ij} = \begin{cases} 1 & \text{if } i \in L, j \in R \text{ and } i \sim j \\ 0 & \text{otherwise} \end{cases}$$

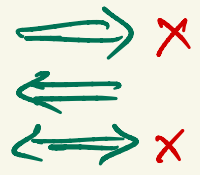


perfect matching:
rook configuration



$$A = (a_{ij})$$

DO $\exists$ perf matching $\leftarrow \leftarrow \det(A) \neq 0$



$$A_G = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix} \mapsto \begin{bmatrix} x_{11} & x_{12} & 0 \\ x_{21} & 0 & x_{23} \\ 0 & 0 & x_{33} \end{bmatrix} = B_G(x_{11} \dots ?)$$



C_4

Thm $\exists$ perf matching $\iff \det B_G \neq 0$

$\iff$ $\checkmark$
 $\implies$ no cancellation over any field

$\det B_G \in F[x_1, \dots, x_m]$ $\deg = n$ unless $-\infty$
works over $\mathbb{F}_p$ where $p \geq 2n$