The (0,1)-Integer Linear Programming Problem

Algorithms – CMSC 37000

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A linear inequality in the variables x, y, z is an inequality of the form $3x - 5y + z \le 6$. (We may replace the coefficients 3, -5, 1, 6 by arbitrary real numbers, and the number of variables may also be arbitrary. We shall assume, however, that all coefficients are integers.) The linear programming (LP) problem takes a system of linear inequalities as input (m inequalities in n variables) and asks its feasibility (does there exist a solution, satisfying all the given inequalities at the same time?). The integer linear programming (ILP) problem asks the existence of a solution in integers (each variable must take an integral value). A (0,1)-ILP asks the existence of a solution where each variable takes the value 0 or 1.

LP is one of the most widely applied algorithmic problems.

Fact. LP is solvable in polynomial time (Khachiyan 1979, Karmarkar 1983).

- (i) Give an example of an LP (with integer coefficients) which is feasible (has solution(s)) but which is not feasible as an ILP. Use as few variables as possible.
- (ii) Same as item (i) but all coefficients in the LP must be 0, 1, or -1, including the numbers on the right hand side.
- (iii) Prove that (0,1)-ILP belongs to NP. State the (polynomial-time verifiable) certificate (of the yes-answers). Indicate why it is verifiable in polynomial time (length of input to verification algorithm, estimated running time of verification).
- (iv) Is it evident that ILP belongs to NP? Argue your answer. Be as specific about the potentially difficult technical detail as you can.
- (v) Assuming 3-COL (3-colorable graphs) is NP-complete, prove that (0,1)-ILP is NP complete, by giving a Karp-reduction from 3-COL to (0,1)-ILP.